

Optimization Methods

Lecture 1

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Material to review: pages 1-15 of Ref[1] and Chapter 1 of Ref [2].

Examples of optimization problems:

- profit or loss in business setting
- speed, distance or fuel consumption in physical problems
- expected return in risky environment
- social welfare problems in the context of government planning

Optimization theory provides
a suitable framework for analyzing and providing solution

Elements of optimization problem: A constraint set X and a cost function f that maps elements of X into a real numbers

- X : The set of constraints of the available decisions x
- $f(x)$: the cost is a scalar measure of undesirability of choosing decision x

Objective: find an optimal decision $x^* \in X$ such that $f(x^*) \leq f(x)$ for $\forall x \in X$

$$\begin{array}{ll} \text{minimize } f(x) & \text{s.t.} \\ x \in X & \end{array} \quad \text{or} \quad \begin{array}{ll} x^* = \operatorname{argmin} f(x) & \text{s.t.} \\ x \in X & \end{array}$$

In our studies $X \subseteq \mathbb{R}^n$ (i.e., $x \in \mathbb{R}^n$)

Different types of optimization: depends on $f(x)$ and X

- **Continuous vs. Discrete**

- Continuous problems: Constraint set X is **infinite** and has “continuous” character

Examples: $X = \mathbb{R}^n$

$$X = \{x \in \mathbb{R}^2 \mid x_2 \geq x_1^2, \quad x_1 + x_2 \leq 2\}$$

Tools to analyze: Mathematics of calculus and convexity

- Discrete problems: mostly because constraint set X is **finite**

Examples: $\left\{ \begin{array}{l} \text{routing} \\ \text{scheduling} \\ \text{Matching} \end{array} \right.$

Important class of discrete problems: integer programming (decision value from some range of integer numbers such as $\{0, 1\}$)

Tools to analyze: Combinational and discrete mathematics; Other special methodology that relate to continuous problems

Different types of optimization: depends on $f(x)$ and X

- **Nonlinear programming**

- cost f is nonlinear and/or
- X is specified by nonlinear equations and inequalities

- **Linear programming**

- cost f is linear
- X is specified by linear inequality constraints

Our focus: nonlinear programming for continuous optimization problems.

$$\begin{aligned} x^* = \operatorname{argmin}_{x \in \mathbb{R}^n} f(x) \quad \text{s.t.} \\ h_i(x) = 0, \quad i \in \{1, \dots, q\} \\ g_i(x) \geq 0, \quad i \in \{1, \dots, p\} \end{aligned}$$

f, h, g : continuously differentiable function of x

e.g., $f \in C^1$ continuously differentiable

e.g., $f \in C^2$ both f and its first derivative are continuously differentiable

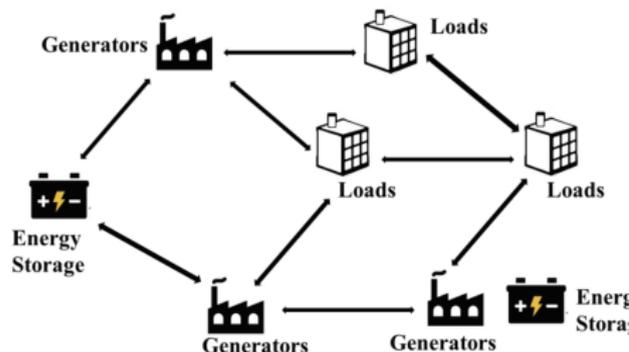
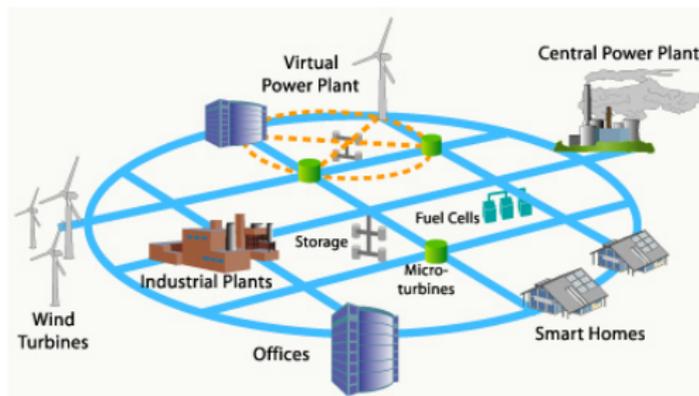
Continuous Optimization: Examples

- Economic Dispatch

$$\text{minimize } f(p) = f^1(p^1) + \dots + f^N(p^N)$$

$$\text{subject to } p_{\min_i} \leq p^i \leq p_{\max_i} \quad i \in \{1, \dots, N\}$$

$$p^1 + \dots + p^N = \text{Demand}$$



- Economic dispatch with storage and transmission loss

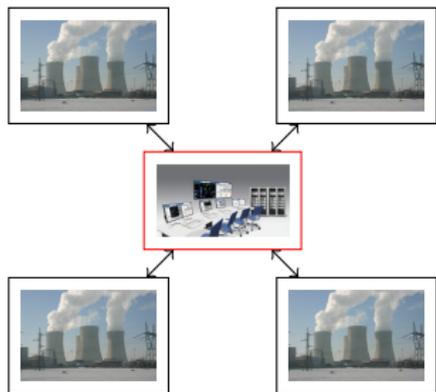
Continuous Optimization: Examples

- Economic Dispatch

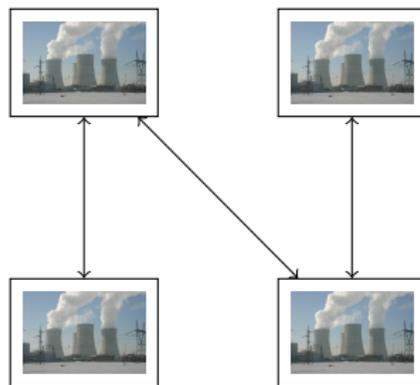
$$\text{minimize } f(p) = f^1(p^1) + \dots + f^N(p^N)$$

$$\text{subject to } p_{\min_i} \leq p^i \leq p_{\max_i} \quad i \in \{1, \dots, N\}$$

$$p^1 + \dots + p^N = \text{Demand}$$



Central Operation



Distributed Operation

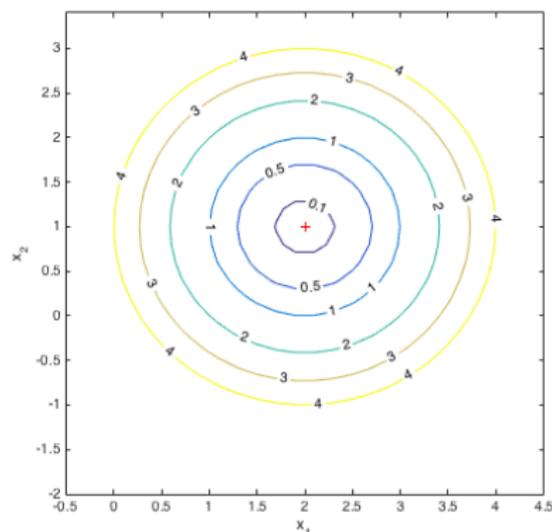
Continuous Optimization: Examples

Unconstraint optimization

$$x^* = \underset{x \in \mathbb{R}^n}{\operatorname{argmin}} f(x)$$

Example:

$$x^* = \underset{x \in \mathbb{R}^2}{\operatorname{argmin}} \underbrace{(x_1 - 2)^2 + (x_2 - 1)^2}_{f(x)}$$



```
x1 = -0.5:0.2:4.5;  
x2 = -2:0.2:3.5;  
[X1,X2] = meshgrid(x1,x2);  
Z = (X1-2).^2+(X2-1).^2;  
v = [0.1,0.1,0.5,0.5,1,1,2,2,3,3,4,4];  
figure  
contour(X1,X2,Z,v, 'ShowText', 'on')  
hold on  
plot(2,1, 'r+')  
pbaspect([1 1 1])  
xlabel('x_1')  
ylabel('x_2')
```

Example:

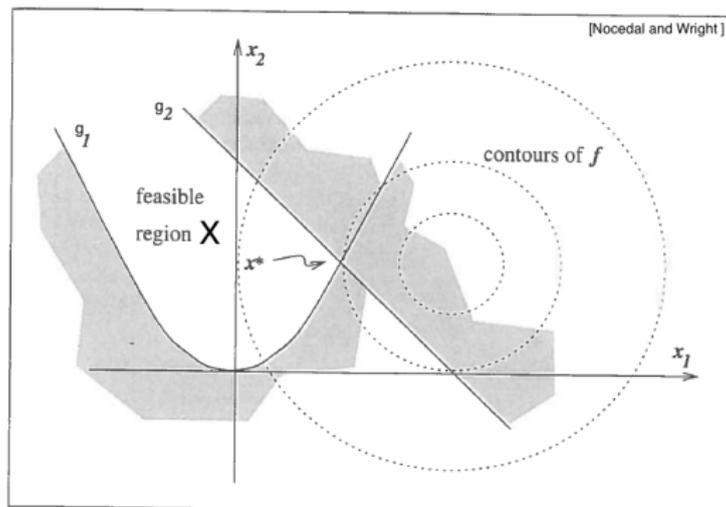
$$x^* = \underset{x \in \mathbb{R}^n}{\operatorname{argmin}} f(x) \quad \text{s.t.}$$

$$h_i(x) = 0, \quad i \in \{1, \dots, q\}$$

$$g_i(x) \geq 0, \quad i \in \{1, \dots, p\}$$

$$x^* = \underset{x \in \mathbb{R}^2}{\operatorname{argmin}} \underbrace{(x_1 - 2)^2 + (x_2 - 1)^2}_{f(x)} \quad \text{s.t.}$$

$$g_i : \begin{cases} -x_1^2 + x_2 \geq 0 \\ -x_1 - x_2 + 2 \geq 0 \end{cases}$$



- $x^* \in \mathbb{R}^n$ is an **unconstrained local minimum** of f if

$$\exists \epsilon > 0 \text{ s.t. } f(x^*) \leq f(x), \quad \forall x \text{ with } \|x - x^*\| < \epsilon,$$

- $x^* \in \mathbb{R}^n$ is an **unconstrained global minimum** of f if

$$f(x^*) \leq f(x), \quad \forall x \in \mathbb{R}^n,$$

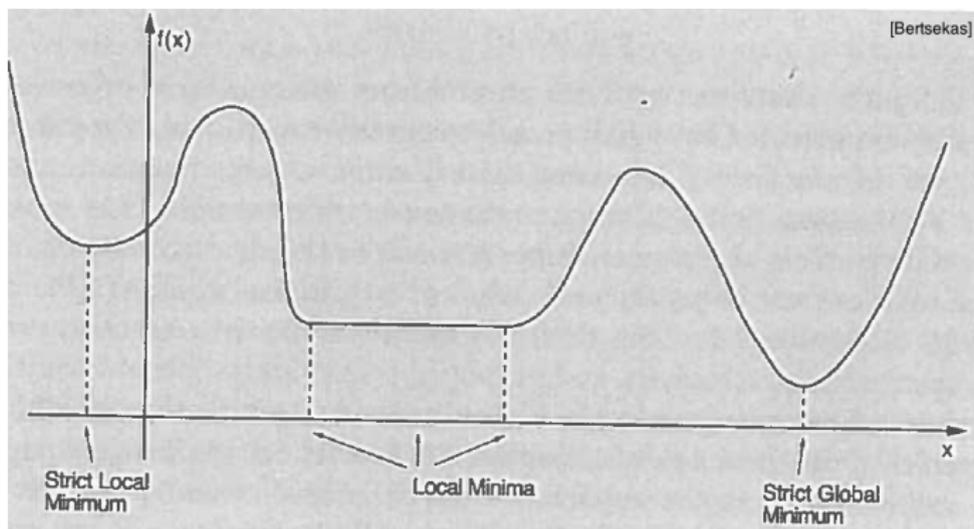
- $x^* \in \mathbb{R}^n$ is an **unconstrained strict local minimum** of f if

$$\exists \epsilon > 0 \text{ s.t. } f(x^*) < f(x), \quad \forall x \text{ with } \|x - x^*\| < \epsilon,$$

- $x^* \in \mathbb{R}^n$ is an **unconstrained strict global minimum** of f if

$$f(x^*) < f(x), \quad \forall x \in \mathbb{R}^n,$$

Local and global minima of a constrained optimization problem



Let X be the constraint set.

- $x^* \in X$ is a **local minimum of f** if

$$\exists \epsilon > 0 \text{ s.t. } f(x^*) \leq f(x), \quad \forall x \in X \text{ with } \|x - x^*\| < \epsilon,$$

- $x^* \in X$ is a **global minimum of f** if

$$f(x^*) \leq f(x), \quad \forall x \in X,$$

- $x^* \in X$ is a **constrained strict local minimum of f** if

$$\exists \epsilon > 0 \text{ s.t. } f(x^*) < f(x), \quad \forall x \in X \text{ with } \|x - x^*\| < \epsilon,$$

- $x^* \in X$ is a **constrained strict global minimum of f** if

$$f(x^*) < f(x), \quad \forall x \in X,$$

- Gradient of a $f(x) : \mathbb{R}^n \rightarrow \mathbb{R}$, $f \in C^1$

$$\nabla f(x) = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \\ \vdots \\ \frac{\partial f}{\partial x_n} \end{bmatrix}$$

Example:

$$f(x) = \frac{1}{2}(x_1 - 2)^2 + x_1 x_2 x_3 + \frac{1}{3}(x_3 + 4)^3$$

$$\nabla f(x) = \begin{bmatrix} (x_1 - 2) + x_2 x_3 \\ x_1 x_3 \\ x_1 x_2 + (x_3 + 4)^2 \end{bmatrix}$$

- Hessian of a $f(x) : \mathbb{R}^n \rightarrow \mathbb{R}$, $f \in C^2$

$$\nabla^2 f(x) = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1 \partial x_1} & \frac{\partial^2 f}{\partial x_1 \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_1 \partial x_n} \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2 \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_2 \partial x_n} \\ \vdots & \vdots & \cdots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1} & \frac{\partial^2 f}{\partial x_n \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_n \partial x_n} \end{bmatrix} = \left[\frac{\partial^2 f}{\partial x_i \partial x_j} \right]_{ij}, \text{ it is a symmetric matrix}$$

Example: $f(x) = \frac{1}{2}(x_1 - 2)^2 + x_1 x_2 x_3 + \frac{1}{3}(x_3 + 4)^3$

$$\nabla^2 f(x) = \begin{bmatrix} 1 & x_3 & x_2 \\ x_3 & 0 & x_1 \\ x_2 & x_1 & 2(x_3 + 4) \end{bmatrix}$$

- [1] Nonlinear Programming: 3rd Edition, by D. P. Bertsekas
- [2] Linear and Nonlinear Programming, by D. G. Luenberger, Y. Ye
- [3] Numerical Optimization, by J. Nocedal and S. J. Wright